

Some questions in typed inquisitive semantics

Lucas Champollion
champollion@nyu.edu

Joint work with Ivano Ciardelli and Floris Roelofsen*

Workshop on questions in logic and semantics
University of Amsterdam
December 15, 2015

Abstract

This talk lays out a compositional account of *wh*-questions in typed inquisitive semantics (Theiler, 2014; Ciardelli and Roelofsen, 2015). Relevant issues include multiple *wh*-questions, the interaction between *wh*-items and disjunction, and *de dicto* readings of *which*-questions.

1 Introduction

- Groenendijk and Stokhof (1984) provided a theory of questions that improved in several respects over Karttunen (1977):
- Basic inquisitive logic (Ciardelli, Groenendijk, and Roelofsen, 2013) in turn improved in some ways on these theories, but did not preserve all of their achievements:

Feature	K77	GS84	InqB
Compositional derivations	yes	yes	no
Interpreting short answers	no	yes	no
De dicto readings of <i>which</i> questions	no	yes	no
Ability to interpret conjoined questions	no	yes	yes
Ability to quantify into questions	no	(yes)	yes
Uniform disjunction across declaratives and interrogatives	no	no	yes
Mention-some readings	no	no	yes
Conditional questions	no	no	yes

- InqB is a logic, and as such does not provide a means to compositionally assign meanings to subsentential constituents. Typed Inquisitive Semantics (Theiler, 2014; Ciardelli and Roelofsen, 2015) provides the bridge between InqB and compositional semantics. We will build on it here.

*Thanks to Theo Janssen for helpful comments. Support from the NYU URCF is gratefully acknowledged.

- As an intermediate step in the compositional derivation, Groenendijk and Stokhof (1984, 1989) compute the *abstract* of a question—an n -place property where n is the number of *wh*-words—and use it to interpret short answers:

- (1) a. Who walks? — John. *Abstract*: $\lambda x.x \text{ walks}$
 b. Who loves whom? — John, Mary. *Abstract*: $\lambda y \lambda x.x \text{ loves } y$

- Typed Inquisitive Semantics gives us the means to compute the abstract of a question.
- Groenendijk and Stokhof (1984) point out that the following inference is invalid when *murderer* is taken *de dicto*:

- (2) a. Holmes knows who is tall.
 b. \Rightarrow Holmes knows which murderer is tall.

- Karttunen (1977) only generates the *de re* reading. The issue has not been revisited in InqB.

2 Typed inquisitive semantics

- Typed inquisitive semantics is a combination of compositional semantics and basic inquisitive logic (Theiler, 2014; Ciardelli and Roelofsen, 2015).
- Possible worlds ($w, w' \dots$) are primitives (type s). States ($p, p' \dots$) are sets of possible worlds (type $\langle s, t \rangle$). Inquisitive propositions ($P, P' \dots$) are sets of states (type $\langle st, t \rangle$).
- We abbreviate $\langle e, \langle et \rangle \rangle$ as $\langle e^2, t \rangle$, and $\langle e, \langle e, \langle et \rangle \rangle \rangle$ as $\langle e^3, t \rangle$, etc. We also write $p(x^n)$ for $p(x_1)(x_2) \dots (x_n)$. Similarly, we write $\lambda x^n.b$ for $\lambda x_1 \lambda x_2 \dots \lambda x_n.b$; and similarly for quantifiers.
- We let *talks* denote the relation that holds between x and p iff p establishes that x talks:

$$(3) \quad \begin{aligned} \llbracket \text{talks} \rrbracket_g &= \lambda x \lambda p \forall w. p(w) \rightarrow \text{talk}(x)(w) \\ &= \lambda x \lambda p. p \subseteq \lambda w. \text{talk}(x)(w) \end{aligned} \quad \text{type } \langle e, \langle st, t \rangle \rangle$$

- We abbreviate $\langle st, t \rangle$ as T . For p_0 a state (type $\langle s, t \rangle$), we write \widehat{p}_0 for the inquisitive proposition $\lambda p. p \subseteq p_0$. Similarly, for p_n of type $\langle e^n, \langle s, t \rangle \rangle$, we write \widehat{p}_n for $\lambda x_1 \dots \lambda x_n \lambda p. p \subseteq \lambda w. p_n(x_1) \dots (x_n)(w)$. For example:

$$(4) \quad \begin{aligned} \llbracket \text{talks} \rrbracket_g &= \widehat{\text{talk}} \\ &= \lambda x \lambda p. p \subseteq \lambda w. \text{talk}(x)(w) \end{aligned} \quad \text{type } \langle e, T \rangle$$

- We represent proper names as constants and use function application to combine meanings:

$$(5) \quad \llbracket \text{John talks} \rrbracket_g = \widehat{\text{talk}}(j) = \lambda p. p \in \widehat{\text{talk}}(j) = \lambda p. p \subseteq \lambda w. \text{talk}(j)(w) \quad \text{type } T$$

3 Propositional connectives

- We assume a type-polymorphic theory of coordination (e.g. Partee and Rooth, 1983). Simplifying slightly, define an inquirable type as either the type T or a type $\langle \alpha, \beta \rangle$ where α is any type and β is an inquirable type.
- We define inquisitive negation, \neg , as in basic inquisitive semantics, and generalize it to higher types:

$$(6) \quad \begin{array}{ll} \text{a.} & \neg_{\langle T, T \rangle} = \lambda P \lambda p. \forall q. P(q) \rightarrow p \cap q = \emptyset & \text{type } \langle T, T \rangle \\ \text{b.} & \neg_{\langle \alpha T, \alpha T \rangle} = \lambda P_{\langle \alpha T \rangle} \lambda x_{\alpha}. \neg_{\langle T, T \rangle} P(x) & \text{type } \langle \alpha T, \alpha T \rangle \end{array}$$

- We represent the meaning of ordinary linguistic negation via inquisitive negation.

$$(7) \quad \llbracket \text{not} \rrbracket_g = \lambda P. \neg P \quad \text{type } \langle \alpha T, \alpha T \rangle$$

For any inquirable type τ we define:

$$(8) \quad \begin{array}{ll} \text{a.} & \llbracket \text{and} \rrbracket_g = \lambda P_{\tau} \lambda Q_{\tau}. P \cap Q & \text{type } \langle \tau, \tau \tau \rangle \\ \text{b.} & \llbracket \text{or} \rrbracket_g = \lambda P_{\tau} \lambda Q_{\tau}. P \cup Q & \text{type } \langle \tau, \tau \tau \rangle \end{array}$$

- As a special case, we will write $\&$ (inquisitive conjunction) for the case where we conjoin two terms P and P' of type T , and similarly for \vee :

$$(9) \quad \begin{array}{ll} \text{a.} & \& \stackrel{\text{def}}{=} \lambda P \lambda P' \lambda p. P'(p) \wedge P(p) \\ \text{b.} & \vee \stackrel{\text{def}}{=} \lambda P \lambda P' \lambda p. P'(p) \vee P(p) \end{array}$$

- Inquisitive conjunction and disjunction share various desirable properties with ordinary conjunction and disjunction, such as idempotence and associativity.
- We assume that proper names can be lifted to generalized quantifiers (note the type):

$$(10) \quad \text{a.} \quad \llbracket \text{Lift(John)} \rrbracket_g = \lambda P_{\langle e, T \rangle}. P(j) \quad \text{type } \langle eT, T \rangle$$

- We can now interpret *John and Mary walk* and *John or Mary walks*.

$$(11) \quad \begin{array}{ll} \text{a.} & \llbracket \text{Lift(John) and Lift(Mary) walk} \rrbracket_g = \widehat{\text{walk}}(j) \& \widehat{\text{walk}}(m) & \text{type } T \\ \text{b.} & \llbracket \text{Lift(John) or Lift(Mary) walks} \rrbracket_g = \widehat{\text{walk}}(j) \vee \widehat{\text{walk}}(m) & \text{type } T \end{array}$$

- *John or Mary walks* is interpreted as an inquisitive proposition with two alternatives:

$$(12) \quad \widehat{\text{walk}}(j) \vee \widehat{\text{walk}}(m) = \lambda p. p \subseteq \lambda w. \text{walk}(j)(w) \vee p \subseteq \lambda w. \text{walk}(m)(w)$$

- We can define type-shifted versions of the inquisitive operators ! (noninquisitive closure) and ? (noninformative closure):

$$(13) \quad ? \stackrel{\text{def}}{=} \lambda P.P \cup \neg P \quad \text{type } \langle \tau, \tau \rangle$$

$$(14) \quad ! \stackrel{\text{def}}{=} \neg \circ \neg \quad \text{type } \langle \tau, \tau \rangle$$

- We assume that any assertion must contain ! at its root.
- This has the following effect (Ciardelli, Groenendijk, and Roelofsen, 2013): Where $A \wp B$ denotes the set of all states that entail A or entail B, $!(A \wp B)$ denotes the set of all states that entail $A \vee B$, including those that do not entail one of the disjuncts.

$$(15) \quad \begin{aligned} & \llbracket ! [\text{John or Mary walk}] \rrbracket_g \\ & = !(\widehat{\text{walk}(j)} \wp \widehat{\text{walk}(m)}) \\ & = \lambda p.p \subseteq \lambda w.\text{walk}(j)(w) \vee \text{walk}(m)(w) \end{aligned} \quad \text{type } T$$

- Finally, we can define inquisitive quantifiers:

$$(16) \quad \begin{aligned} \text{a.} \quad & \exists x\phi \stackrel{\text{def}}{=} \lambda p \exists x \phi(p) \\ \text{b.} \quad & \forall x\phi \stackrel{\text{def}}{=} \lambda p \forall x \phi(p) \end{aligned}$$

4 Wh-questions and the abstract

- We assume that questions, whether embedded or not, are headed by a silent Q morpheme (Baker, 1970), which projects an *interrogative nucleus*. The complement of Q is the *abstract*.

$$(17) \quad \begin{array}{c} \text{interrogative} \\ \text{nucleus} \\ \diagup \quad \diagdown \\ Q \quad \text{abstract} \end{array}$$

- The abstract is of type $\langle e^n, T \rangle$: e.g. $\langle e, T \rangle$ for single-*wh* questions, $\langle e, eT \rangle$ for double-*wh*-questions.
- We could naively assume that *wh*-phrases like *who* are identity functions:

$$(18) \quad \text{e.g. } \llbracket \text{who} \rrbracket_g \text{ in subject position} = \lambda P_{et} \lambda x_e.P(x)$$

- This differs from the treatment in Theiler (2014), where *wh*-phrases are treated as inquisitive existentials.
- But this will not work when we need to pass the abstract across sentence boundaries:

$$(19) \quad \text{Whom do you want Mary to invite?}$$

- Here, *want* expects a proposition, so *whom* must leave a trace behind or be interpreted in situ at LF.

- This process can violate islands, so an in-situ based account is preferable (cf. Reinhart, 1997):

- (20) a. Who thinks that who walks?
b. Who will be offended if we invite whom?

- So we assume instead, following Baker (1970), that *who* carries an index, that the Q morpheme binds such indices or triggers lambda abstraction below it:

- (21) a. Who walks? $\sim \rightarrow [Q [1 [who_1 \text{ walks}]]]$
b. Who loves whom? $\sim \rightarrow [Q [1 [2 [who_1 \text{ loves whom}_2]]]]$
c. Who thinks that who walks? $\sim \rightarrow [Q [1 [2 [who_1 \text{ thinks [that who}_2 \text{ walks}]]]]]]$

- Sometimes the abstract will be noninquisitive:

- (22) a. $[[[1 [who_1 \text{ walks}]]]] = \lambda x_1. \widehat{\text{walk}}(x_1)$
b. $[[[1 [2 [who_1 \text{ loves whom}_2]]]]] = \lambda x_1 \lambda x_2. \widehat{\text{love}}(x_2)(x_1)$

- Sometimes it will be inquisitive:

- (23) Who walks or talks?
a. $[Q [1 [who_1 [\text{walks or talks}]]]]$
b. $[[[1 [who_1 [\text{walks or talks}]]]]] = \lambda x_1. \widehat{\text{walk}}(x_1) \vee \widehat{\text{talk}}(x_1)$

- The basic meaning InqB assigns to (21a) and (21b) captures their mention-some reading:

- (24) a. $? \exists x. \widehat{\text{walk}}(x)$
b. $? \exists x \exists y. \widehat{\text{love}}(y)(x)$

- For example, (24a) has the following alternatives:

that John walks, that Mary walks, ..., that nobody walks

- It would be a mistake to treat inquisitive abstracts in the same way, however:

- (25) $? \exists x. \widehat{\text{walk}}(x) \vee \widehat{\text{talk}}(x)$

- This has the following alternatives: *that John walks, that John talks, that Mary walks, that Mary talks, ..., and that nobody walks or talks.*

- A better translation uses noninquisitive closure:

- (26) $? \exists x. !(\widehat{\text{walk}}(x) \vee \widehat{\text{talk}}(x))$

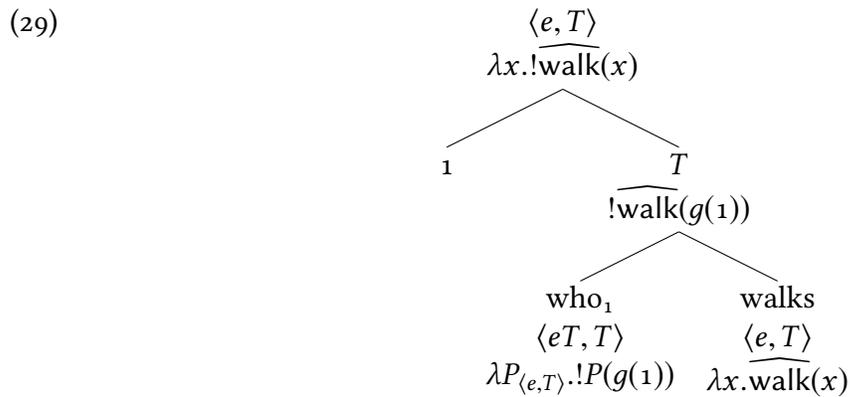
- This has the alternatives *that John walks or talks, that Mary walks or talks, ..., that nobody walks or talks.*

- What is responsible for the introduction of noninquisitive closure?
- Compositionally, we seem to have two options: ! is introduced by Q , or by the wh -phrases.
- In non- wh questions, Q often does not seem to introduce !

(27) Would you like coffee \uparrow , or tea \downarrow ?
 \neq Is it the case that you would like either coffee or tea?

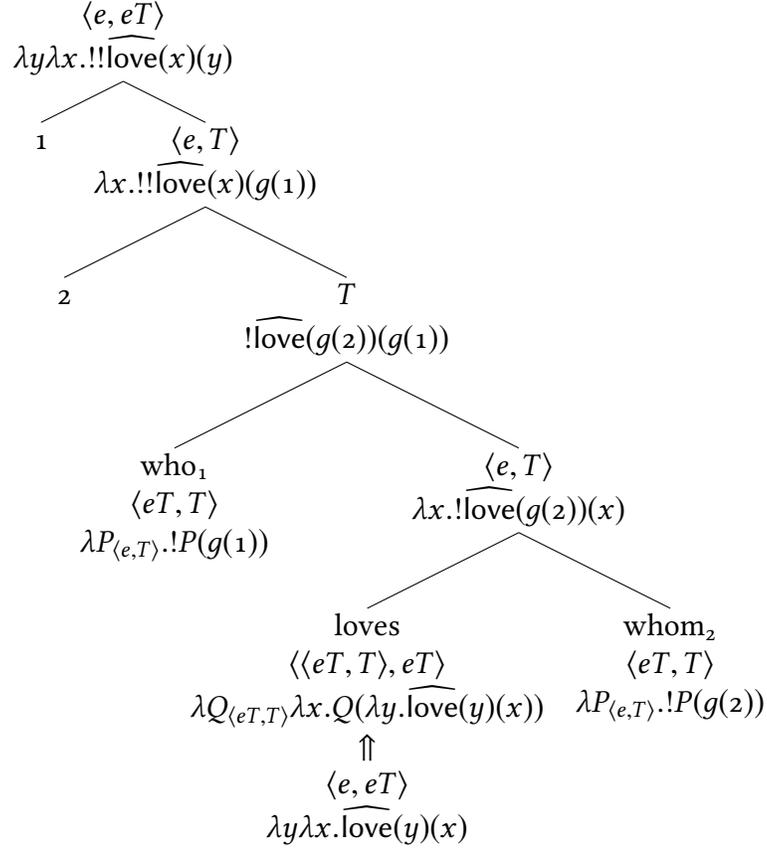
- So we assume that it is the wh -phrases that are responsible for the introduction of !.

(28) $[[\text{who}_i]]_g = \lambda P_{\langle e, T \rangle} !P(g(i))$ type $\langle eT, T \rangle$



- In nonsubject position, we resolve type mismatches by type-shifters on the verb (Hendriks, 1993).

(30)



5 The Q operator

- The Q operator maps abstracts to inquisitive propositions.
- As is well known, there are several relevant candidate propositions:

(31) John knows who is tall.

- | | | |
|----|--|----------------------------|
| a. | John knows of some x that x is tall. | <i>mention-some</i> |
| b. | John knows of every tall x that x is tall. | <i>weakly exhaustive</i> |
| c. | John knows of every x whether x is tall. | <i>strongly exhaustive</i> |

- Groenendijk and Stokhof (1984) take (31c) as basic, which makes it hard to model (31a) and (31b) (Heim, 1994; Beck and Rullmann, 1999).
- InqB takes (31a) as basic, so one can model (31b) and (31c) through additional operations (Theiler, 2014).
- We assume that exhaustification optionally takes place within the interrogative nucleus; the precise “flavor” of exhaustivity is determined higher up (Theiler, 2014).
- We base the meaning of Q on the inquisitive existential \exists and on the operator $?$. (This is a simplification. For certain purposes involving non-*wh* questions, it would be more

accurate to use $\langle ? \rangle$, which leaves inquisitive meanings alone, and applies $?$ to noninquisitive meanings.)

$$(32) \quad \llbracket Q^n \rrbracket_g = \lambda P_{\langle e^n, T \rangle}. ? \exists x^n. P(x^n) \quad \text{type } \langle e^n T, T \rangle$$

- Some special cases:

$$(33) \quad \begin{array}{ll} \text{a. } \llbracket Q^0 \rrbracket_g \text{ (for non-} wh \text{ questions)} = \lambda P_T. ? P & \text{type } \langle T, T \rangle \\ \text{b. } \llbracket Q^1 \rrbracket_g \text{ (for single-} wh \text{ questions)} = \lambda P_{\langle e, T \rangle}. ? \exists x. P(x) & \text{type } \langle eT, T \rangle \\ \text{c. } \llbracket Q^2 \rrbracket_g \text{ (for double-} wh \text{ questions)} = \lambda R_{\langle e, eT \rangle}. ? \exists x \exists y. R(x)(y) & \text{type } \langle \langle e, eT \rangle, T \rangle \end{array}$$

- A second version of the operator has exhaustivity built in:

$$(34) \quad \llbracket Q_{+exh}^n \rrbracket_g = \lambda R_{\langle e^n, T \rangle}. \lambda p. \forall q \subseteq p. (\lambda x^n. R(x^n)(q) = \lambda x^n. R(x^n)(p)) \quad \text{type } \langle e^n T, T \rangle$$

- Some special cases:

$$(35) \quad \begin{array}{ll} \text{a. } \llbracket Q_{+exh}^0 \rrbracket_g \text{ (for non-} wh \text{ questions)} \\ = \lambda P_T. ? P = \lambda P_T. P \vee \neg P & \text{type } \langle T, T \rangle \\ \text{b. } \llbracket Q_{+exh}^1 \rrbracket_g \text{ (for single-} wh \text{ questions)} \\ = \lambda P_{\langle e, T \rangle}. \lambda p. \forall q \subseteq p. (\lambda x. P(x)(q) = \lambda x. P(x)(p)) & \text{type } \langle eT, T \rangle \\ \text{c. } \llbracket Q_{+exh}^2 \rrbracket_g \text{ (for double-} wh \text{)} \\ = \lambda R_{\langle e, eT \rangle}. \lambda p. \forall q \subseteq p. (\lambda y \lambda x. R(y)(x)(q) = \lambda y \lambda x. R(y)(x)(p)) & \text{type } \langle \langle e, eT \rangle, T \rangle \end{array}$$

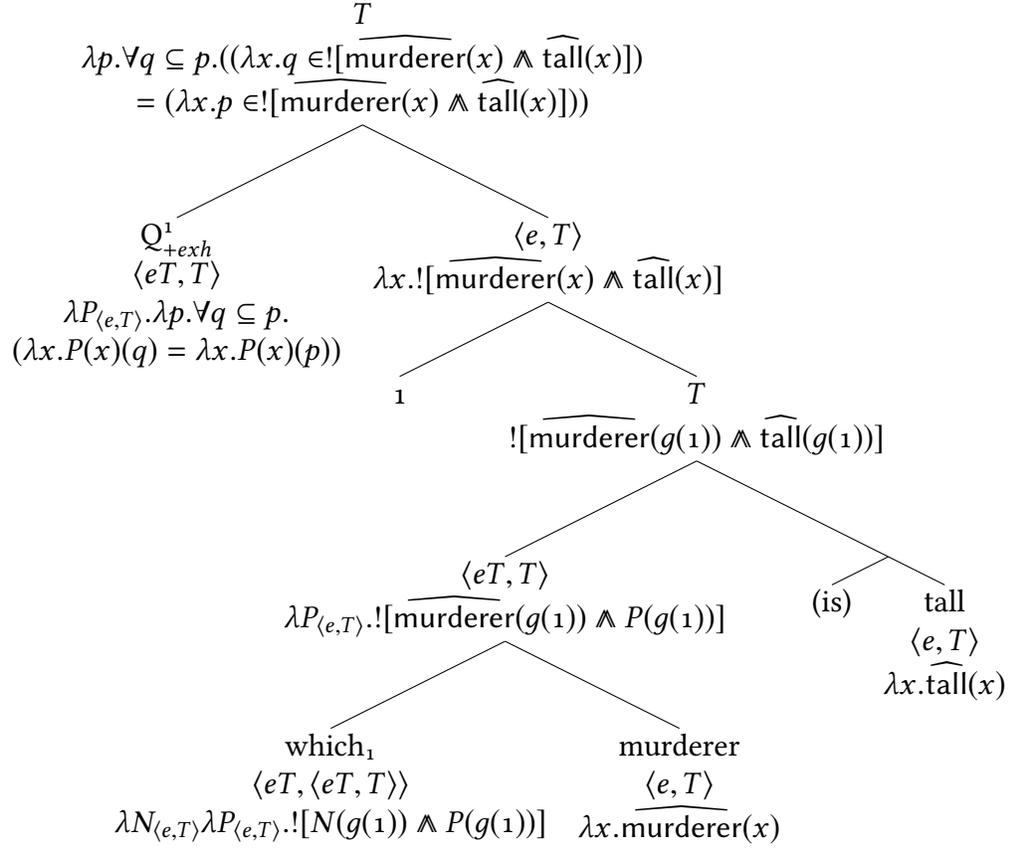
6 Which-questions

- Following Groenendijk and Stokhof (1984), we assume that *which*-questions interpret their noun *in situ*. (We give here a non-presuppositional account but we are working on a presuppositional extension.)

$$(36) \quad \llbracket \text{which}_i \rrbracket_g = \lambda N_{\langle e, T \rangle} \lambda P_{\langle e, T \rangle}. ! [N(g(i)) \text{ \# } P(g(i))] \quad \text{type } \langle eT, \langle eT, T \rangle \rangle$$

- As for *who*, in nonsubject position we resolve type mismatches by type-shifters on the verb.

(37)



- This gives us access to the kind of object we need in order to compute a *de dicto* reading.
- Following Groenendijk and Stokhof (1984) we have given what amounts to a symmetric account. Of course, we know this can't be the whole story:

(38) From Higginbotham (1996):

- Which men are bachelors?
- #Which bachelors are men?

- We are currently studying presuppositional extensions of inquisitive semantics that would allow us to import accounts such as Rullmann and Beck (1998) that capture the contrast between these questions in terms of the presuppositions of the *which*-phrase.

References

- Baker, Carl Lee (1970). Notes on the description of English questions: the role of an abstract question morpheme. *Foundations of Language* 6(2), 197–219. URL: <http://www.jstor.org/stable/25000451>.
- Beck, Sigrid and Hotze Rullmann (1999). A flexible approach to exhaustivity in questions. *Natural Language Semantics* 7(3), 249–298. DOI: 10.1023/a:1008373224343.

- Ciardelli, Ivano, Jeroen Groenendijk, and Roelofsen (2013). Inquisitive semantics: A new notion of meaning. *Language and Linguistics Compass* 7(9), 459–476. DOI: 10.1111/lnc3.12037.
- Ciardelli, Ivano and Floris Roelofsen (2015). Alternatives in Montague grammar. In: *Proceedings of Sinn und Bedeutung 20*. Ed. by Eva Csipak and Hedde Zeijlstra. Göttingen, Germany: University of Göttingen, pp. 161–178. URL: <http://semanticsarchive.net/Archive/jQzODVh0/>.
- Groenendijk, Jeroen and Martin Stokhof (1984). Studies on the semantics of questions and the pragmatics of answers. PhD thesis. Amsterdam, Netherlands: University of Amsterdam.
- Groenendijk, Jeroen and Martin Stokhof (1989). Type-shifting rules and the semantics of interrogatives. In: *Properties, types and meaning. Volume II: semantic issues*. Ed. by Gennaro Chierchia, Barbara Hall Partee, and Raymond Turner. Vol. 39. Studies in Linguistics and Philosophy. Amsterdam, Netherlands: Springer, pp. 21–68. DOI: 10.1007/978-94-009-2723-0_2.
- Heim, Irene (1994). Interrogative semantics and Karttunen’s semantics for *know*. In: *The Israeli Association for Computational Linguistics: Proceedings of the 9th Annual Conference and Workshop on Discourse*. Ed. by Rhonna Buchalla and Anita Mittwoch. Vol. 1. Jerusalem, Israel: Akademon, pp. 128–144.
- Hendriks, Herman (1993). Studied flexibility. PhD thesis. Amsterdam, Netherlands: University of Amsterdam.
- Higginbotham, James (1996). The semantics of questions. In: *Handbook of contemporary semantic theory*. Ed. by Shalom Lappin. Basil Blackwell. DOI: 10.1111/b.9780631207498.1997.00017.x.
- Karttunen, Lauri (1977). Syntax and semantics of questions. *Linguistics and Philosophy* 1(1), 3–44. DOI: 10.1007/bf00351935.
- Partee, Barbara H. and Mats Rooth (1983). Generalized conjunction and type ambiguity. In: *Meaning, use and interpretation of language*. Ed. by Rainer Bäuerle, Christoph Schwarze, and Arnim von Stechow. Berlin, Germany: de Gruyter, pp. 361–383. DOI: 10.1515/9783110852820.361.
- Reinhart, Tanya (1997). Quantifier scope: How labor is divided between QR and choice functions. *Linguistics and Philosophy* 20(4), 335–397. DOI: 10.1023/a:1005349801431.
- Rullmann, Hotze and Sigrid Beck (1998). Presupposition projection and the interpretation of *which* questions. In: *Proceedings of the 8th Semantics and Linguistic Theory Conference (SALT 8)*. Ed. by Devon Strolovitch and Aaron Lawson. Ithaca, NY: Cornell University, pp. 215–232. DOI: 10.3765/salt.v8i0.2811.
- Theiler, Nadine (2014). A multitude of answers: embedded questions in typed inquisitive semantics. MA thesis. University of Amsterdam.